



**Chapter 8** 

### Subroutine

#### **Expected Outcomes**

- Describe and apply the stack for data storage
- Describe the process of subroutine in any programs
- Develop a subroutine and code
- Interpret subroutine process in the stack
- Write and calculate a delay subroutine



# **Stack**

- Stack is a special area in memory and normally it is used to keep track of and store CPU register information during execution
- Most stack uses LIFO concept
- It has a stack pointer (SP) to indicate where to push or pull data and A7 is used for this purpose
- To store data, CPU pushes it onto the stack and then decrements the SP
- To recover data from the stack, it increments the SP and then pulls the data
- The stack grows toward low memory addresses

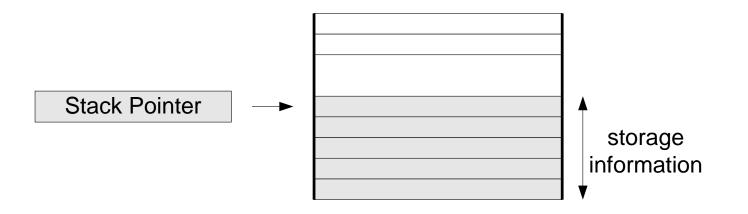


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# **Stack Pointer**

Stack must be located in RAM location and normally place above program and data

**Stack pointer** must be initialized at the beginning of the program





# **Push & Pull Instruction**



 ARI with pre-decrement and ARI with post-increment are required to perform the push and pull operation
Push:

MOVE.s Source,-(SP)



MOVE.s (SP)+, destination



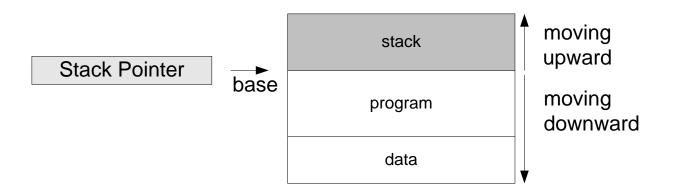


# **Initialize Stack Pointer**

To initialize stack pointer

MOVEA.L #BASE, SP

where BASE is the end address of a stack







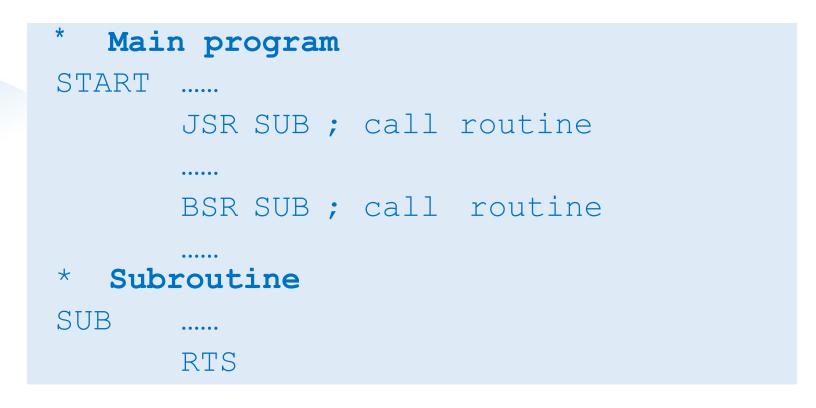
# **Subroutine**

- Subroutine is a section of a program that may be used one or more times
- With subroutine, the program is much simpler, short, efficient and more understandable
- The main program calls subroutine to perform certain steps using the instruction JSR (jump to subroutine) or BSR (branch to subroutine)
- It executes the subroutines until the instruction RTS (return from subroutine)
- It returns to main program and continues at the instruction following instruction JSR or BSR





# **Subroutine**

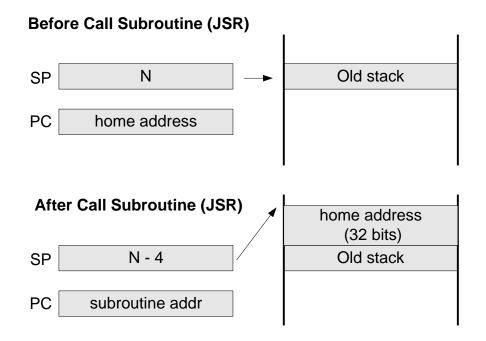




# **JSR Operation**



In order to return to main program, the current PC must be stored in the stack



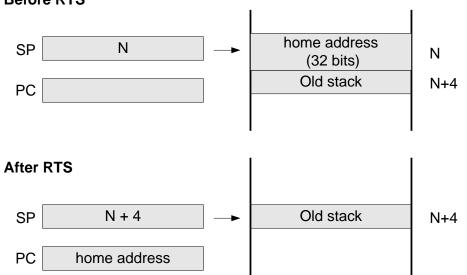




# **RTS Operation**

In order to return to main program, the return address of main program must be placed back to the PC.

**RTS will ensure the procedure is followed** 



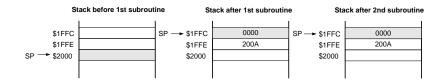




### **JSR&RTS Operation**



		ORG \$2000
	STACK	EQU *
2000	START	MOVEA #STACK,
2004		NOP
2006		JSR SUB
200A	RET1	NOP
200C		JSR SUB
2010	RET2	NOP
2012		BRA *
2014	SUB	NOP
2016		RTS





SP



# **Protecting Registers**

- When subroutine is executed, the content of registers may alter if the registers are used in the subroutine
- Thus, the register need to be stored in stack
- Following is one way to secure the content of registers (In this case, register D1-D2 and A3, A6)

MOVE.L	D1,	-(SP)		
MOVE.L	D2,	-(SP)		
MOVE.L	A3,	-(SP)		
MOVE.L	A6,	-(SP)		
content of subroutine				
MOVE.L	(SP)	+, A6		
MOVE.L	(SP)	+, A3		
MOVE.L	(SP)	+, D2		
MOVE.L	(SP)	+, D1		
RTS				



### MOVEM Instruction



Another alternative way to store the content of register is using the MOVEM (move multiple registers) instruction

■ Syntax

MOVEM.s <list registers>,-SP

MOVEM.s SP+, <list registers>

Following is an example to store register D1-D3 and A3, A4, A6 MOVEM.L D1-D3/A3-A4/A6, -(SP) ...content of subroutine... MOVEM.L (SP)+, D1-D3/A3-A4/A6 RTS



# **Macros vs Subroutines**



- Both permits a group of instruction to be defined in a single entity with a unique given label or name called up when needed
- A subroutine is called by **BSR** or **JSR** instructions, while macro is called by simply its name
- Macros are not substitute for subroutines
- Support for subroutines is provided by CPU as it is part of instruction set, while support for macros is part of the assembler





# **Self-Test**

#### Exercise

If SP=\$0040000C and PC=\$00400500, what is the value of SP when JSR \$00400600 is executed ?

#### **Exercise**:

Calculate the value of SP if the following program is executed

ORG \$4000 START MOVEA #\$2000,SP MOVEM.L D0-D2/A0/A4-A6,-(SP)

